

Post-Quantum Signatures from Secure Multiparty Computation

Thibault Feneuil

Journées C2 - 2023

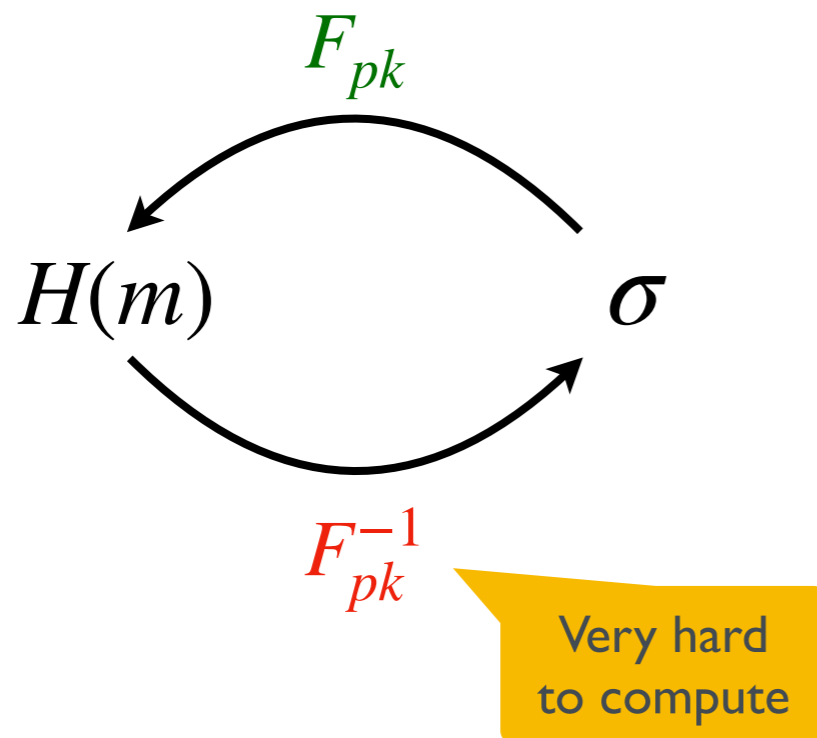
October 16, 2023, Najac (France)



Introduction

How to build signature schemes?

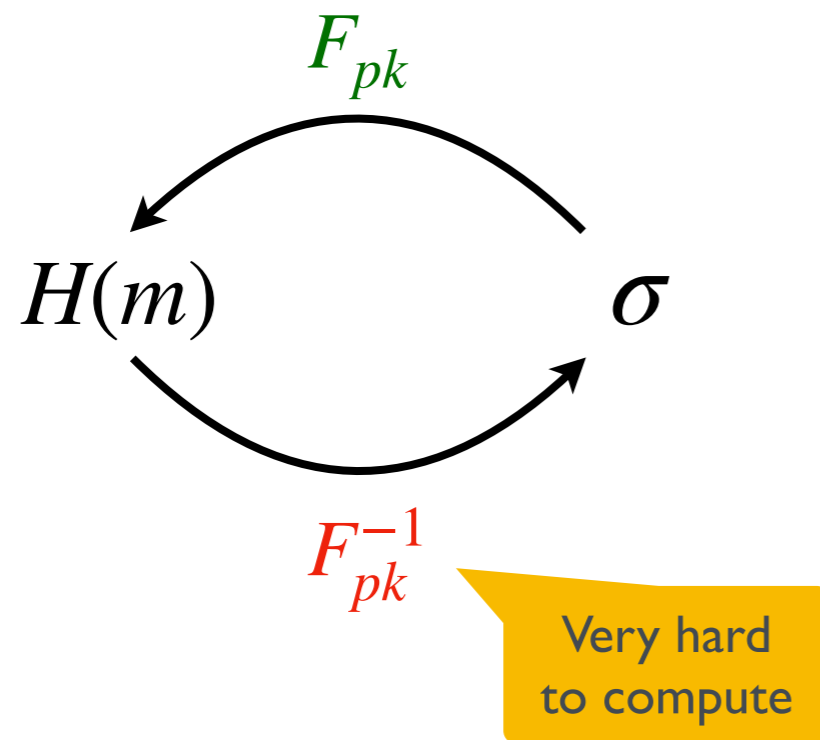
Hash & Sign



- Short signatures
- “Trapdoor” in the public key

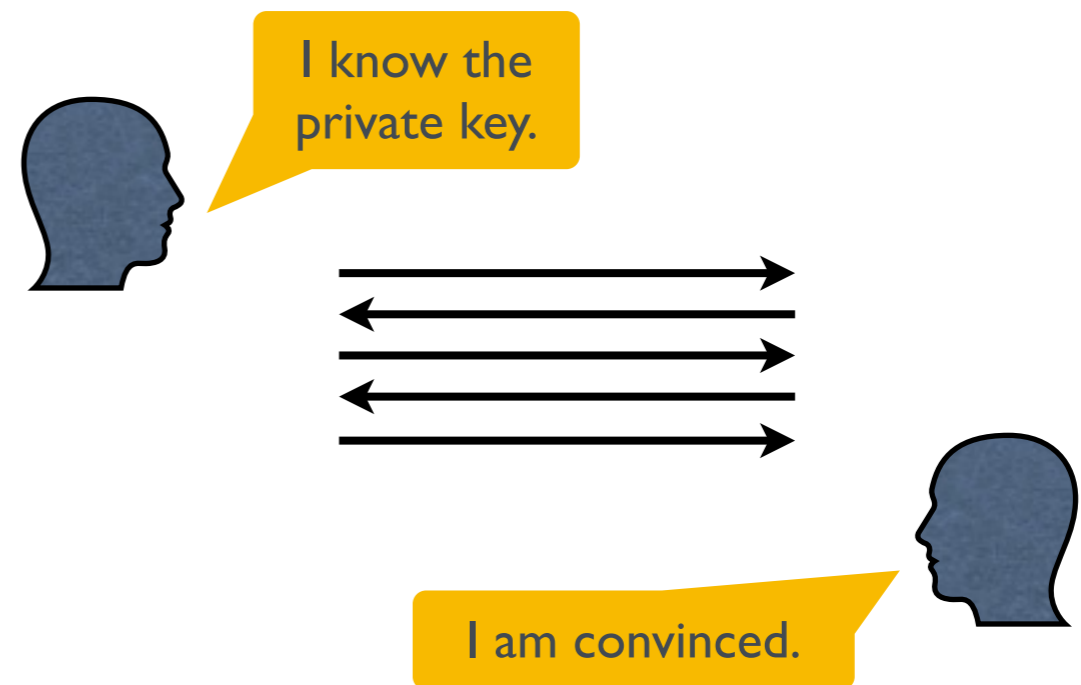
How to build signature schemes?

Hash & Sign



- Short signatures
- “Trapdoor” in the public key

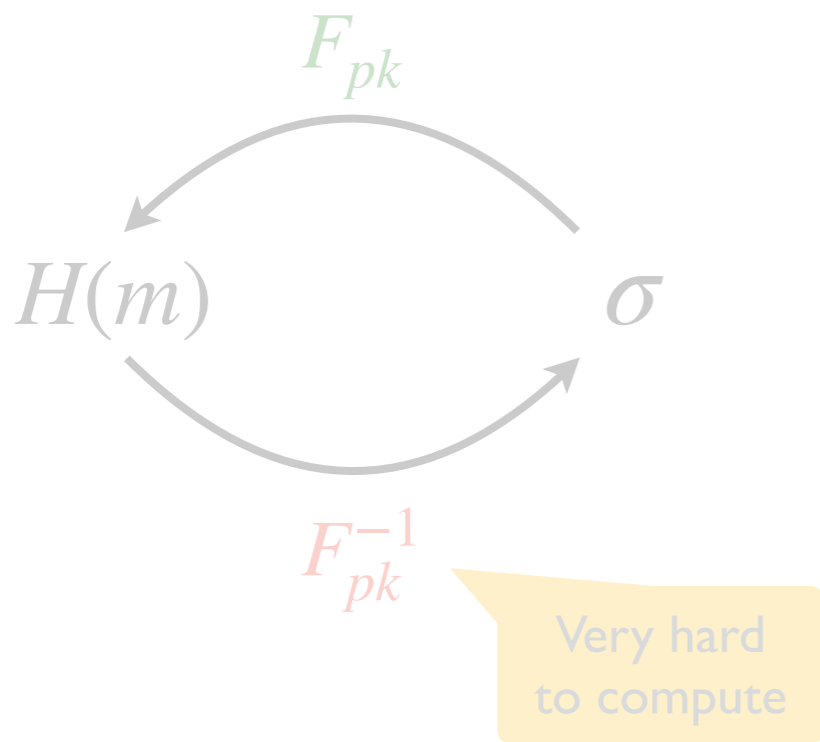
From a zero-knowledge proof



- Large(r) signatures
- Short public key

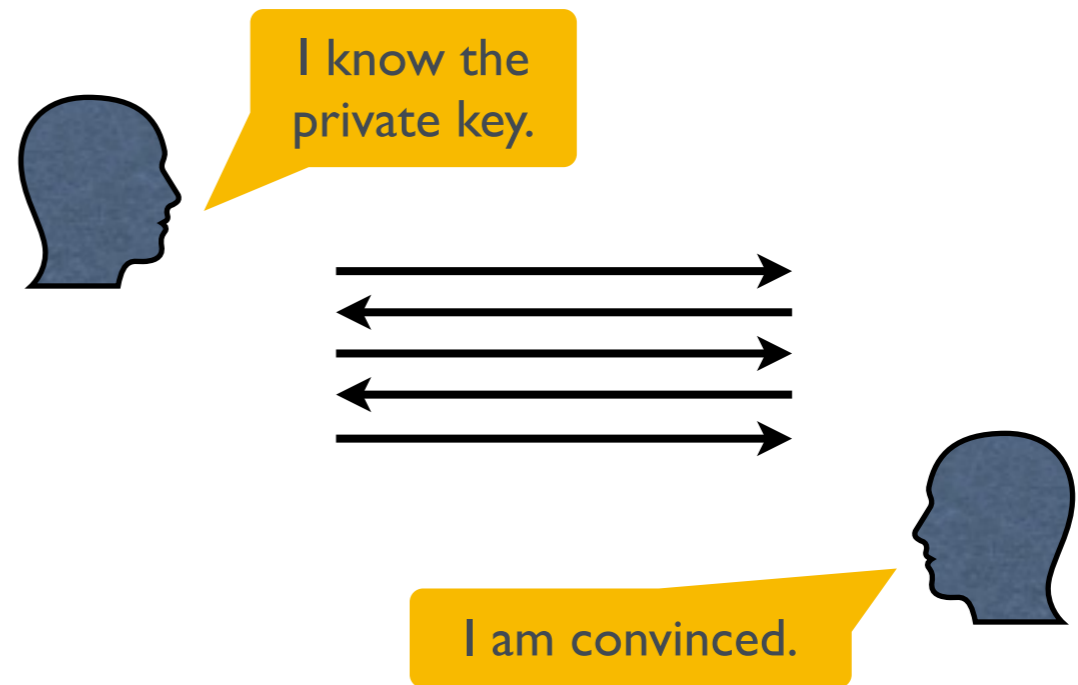
How to build signature schemes?

Hash & Sign



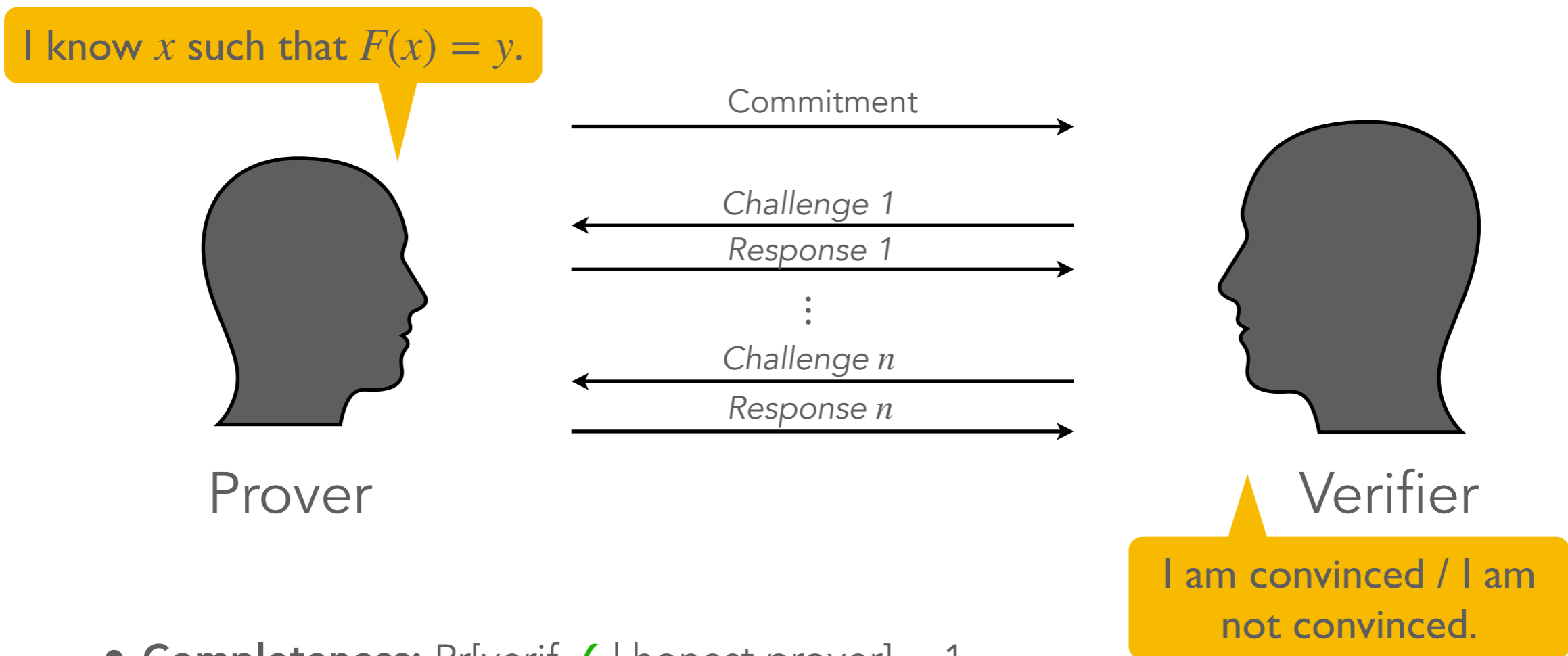
- Short signatures
- “Trapdoor” in the public key

From a zero-knowledge proof



- Large(r) signatures
- Short public key

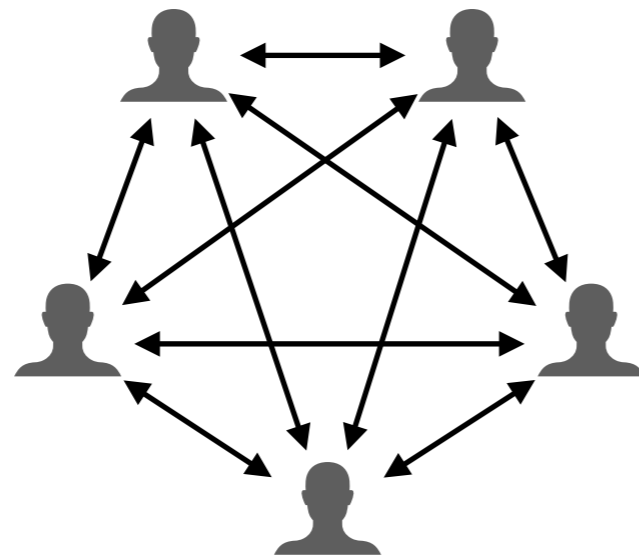
Proof of knowledge



- **Completeness:** $\Pr[\text{verif } \checkmark \mid \text{honest prover}] = 1$
- **Soundness:** $\Pr[\text{verif } \checkmark \mid \text{malicious prover}] \leq \varepsilon$ (e.g. 2^{-128})
- **Zero-knowledge:** verifier learns nothing on x

MPC in the Head

- [IKOS07] Yuval Ishai, Eyal Kushilevitz, Rafail Ostrovsky, Amit Sahai: “Zero-knowledge from secure multiparty computation” (STOC 2007)
- Turn a *multiparty computation* (MPC) into an identification scheme / zero knowledge proof of knowledge



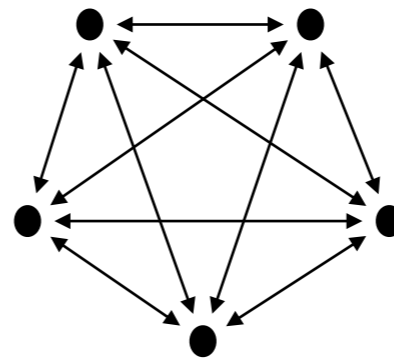
- **Generic:** can be apply to any cryptographic problem

One-way function

$$F : x \mapsto y$$

E.g. AES, MQ system,
Syndrome decoding

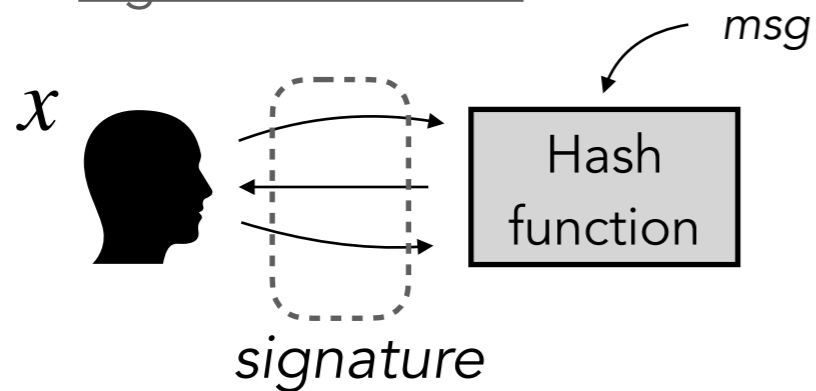
Multiparty computation (MPC)



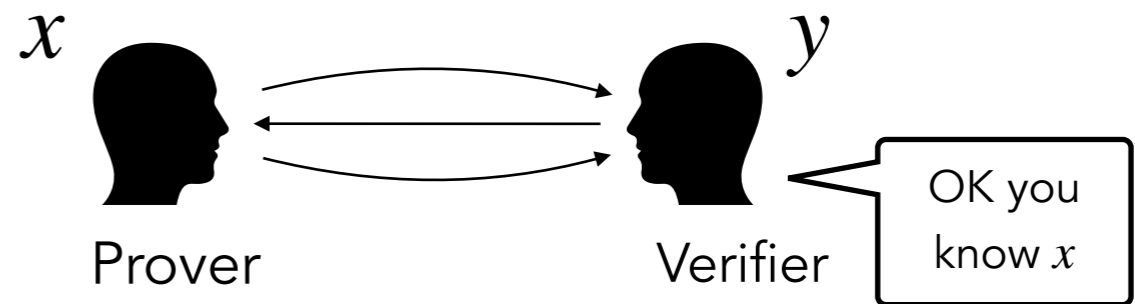
Input sharing $[[x]]$
Joint evaluation of:

$$g(x) = \begin{cases} \text{Accept} & \text{if } F(x) = y \\ \text{Reject} & \text{if } F(x) \neq y \end{cases}$$

Signature scheme



Zero-knowledge proof

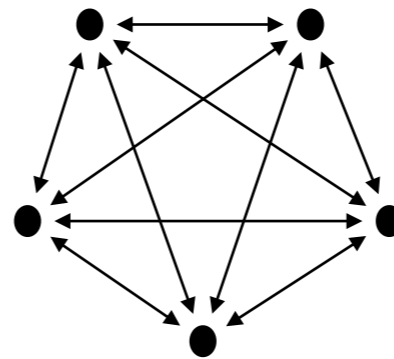


One-way function

$$F : x \mapsto y$$

E.g. AES, MQ system,
Syndrome decoding

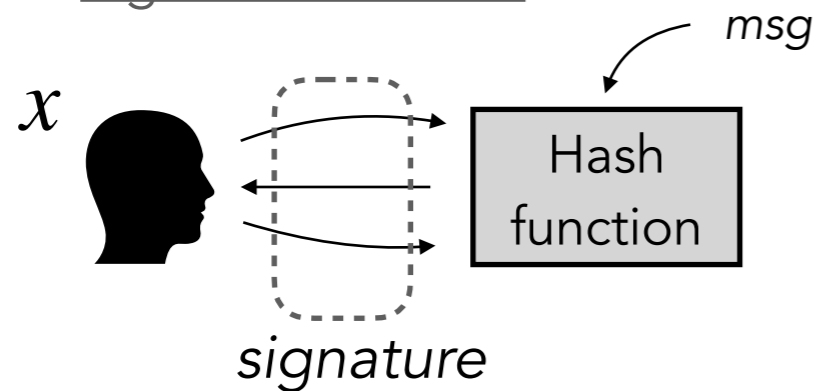
Multiparty computation (MPC)



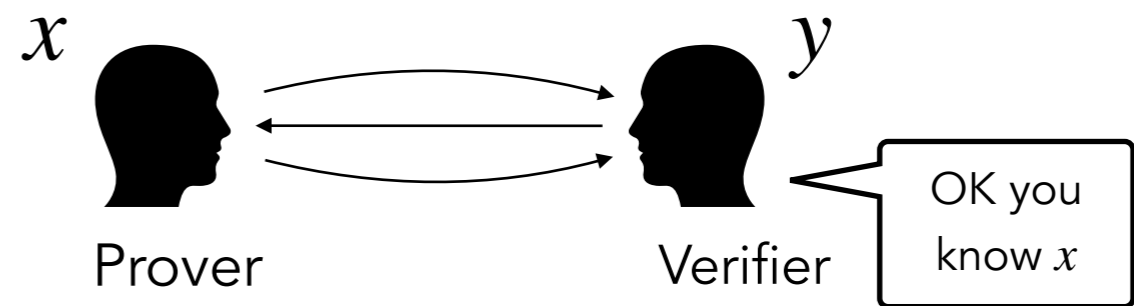
Input sharing $[[x]]$
Joint evaluation of:

$$g(x) = \begin{cases} \text{Accept} & \text{if } F(x) = y \\ \text{Reject} & \text{if } F(x) \neq y \end{cases}$$

Signature scheme



Zero-knowledge proof



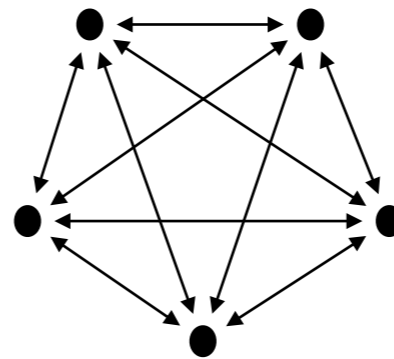
$$[[x]] = ([[x]]_1, \dots, [[x]]_N) \quad \text{s.t.} \quad x = [[x]]_1 + \dots + [[x]]_N$$

One-way function

$$F : x \mapsto y$$

E.g. AES, MQ system,
Syndrome decoding

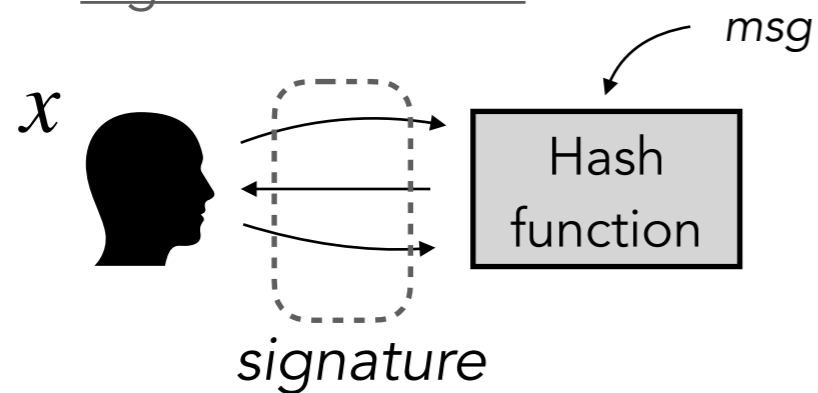
Multiparty computation (MPC)



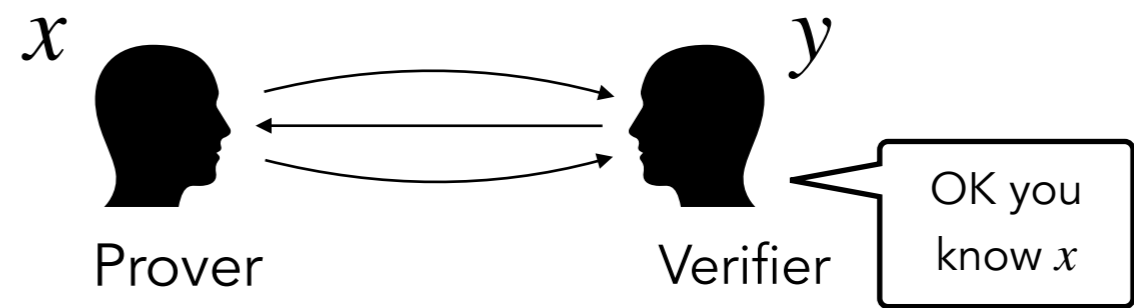
Input sharing $[[x]]$
Joint evaluation of:

$$g(x) = \begin{cases} \text{Accept} & \text{if } F(x) = y \\ \text{Reject} & \text{if } F(x) \neq y \end{cases}$$

Signature scheme



Zero-knowledge proof

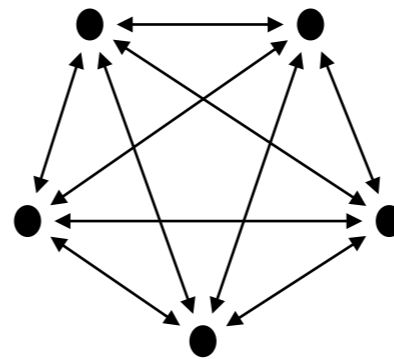


One-way function

$$F : x \mapsto y$$

E.g. AES, MQ system,
Syndrome decoding

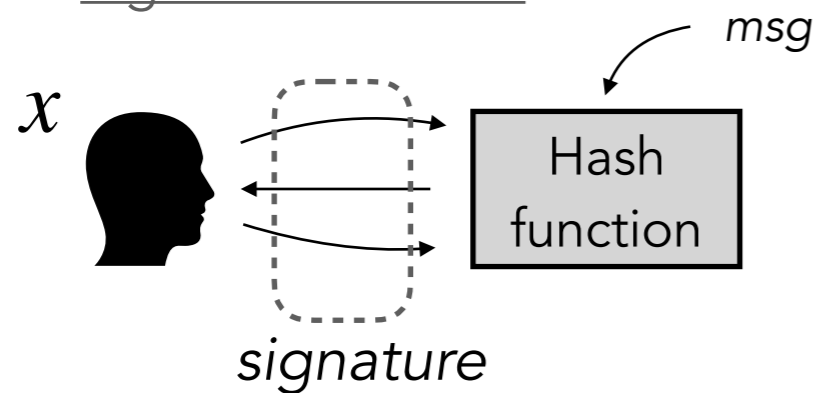
Multiparty computation (MPC)



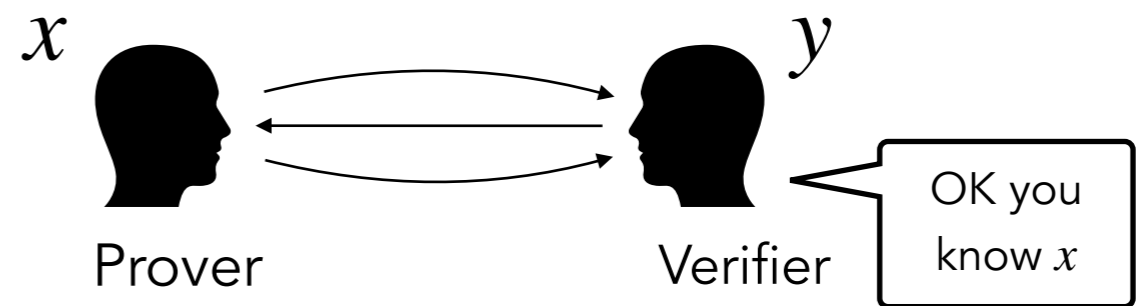
Input sharing $[[x]]$
Joint evaluation of:

$$g(x) = \begin{cases} \text{Accept} & \text{if } F(x) = y \\ \text{Reject} & \text{if } F(x) \neq y \end{cases}$$

Signature scheme



Zero-knowledge proof

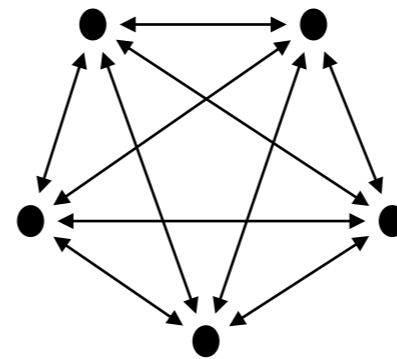


One-way function

$$F : x \mapsto y$$

E.g. AES, MQ system,
Syndrome decoding

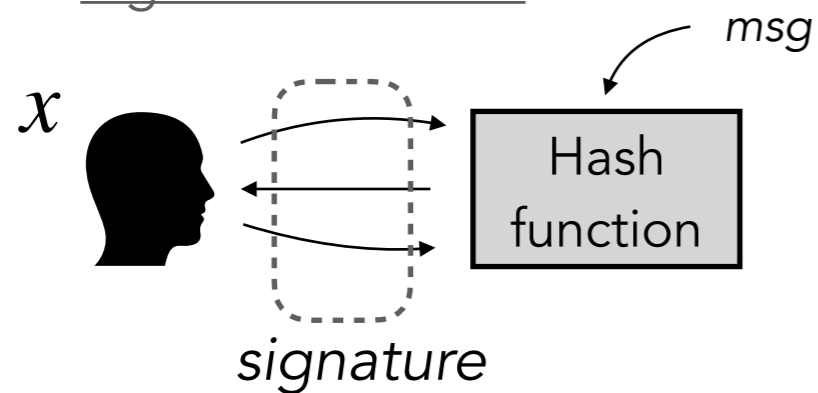
Multiparty computation (MPC)



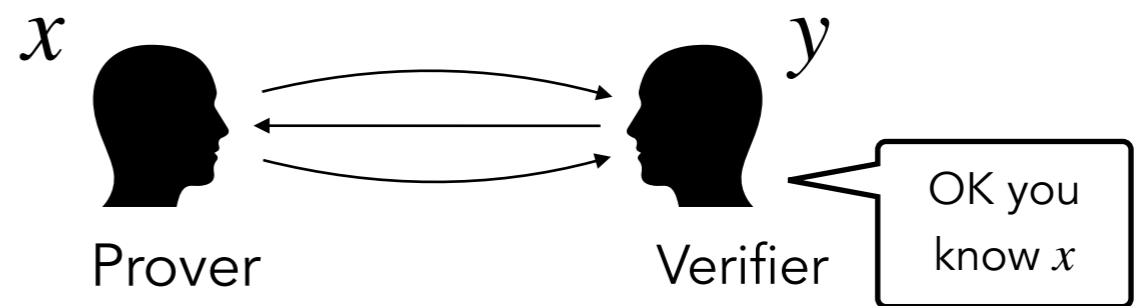
Input sharing $[[x]]$
Joint evaluation of:

$$g(x) = \begin{cases} \text{Accept} & \text{if } F(x) = y \\ \text{Reject} & \text{if } F(x) \neq y \end{cases}$$

Signature scheme



Zero-knowledge proof

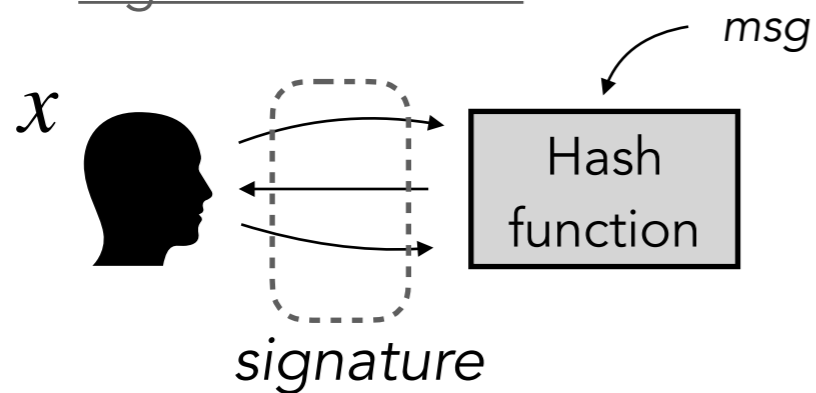


One-way function

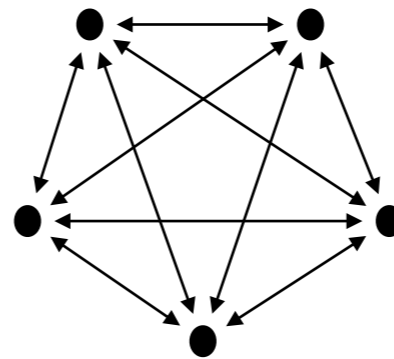
$$F : x \mapsto y$$

E.g. AES, MQ system,
Syndrome decoding

Signature scheme



Multiparty computation (MPC)

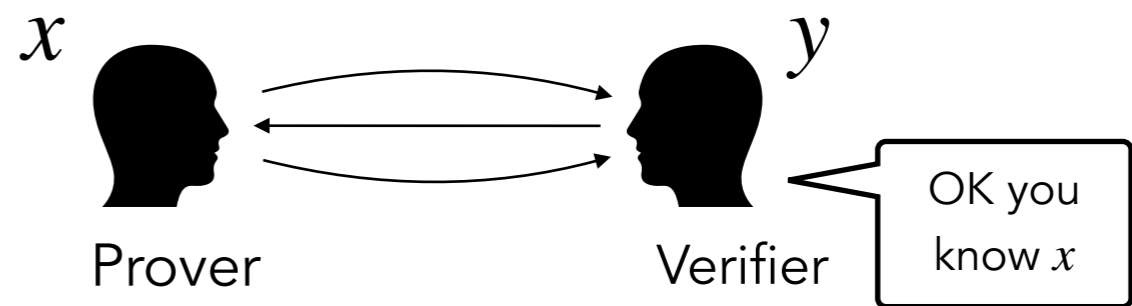


Input sharing $[[x]]$
Joint evaluation of:

$$g(x) = \begin{cases} \text{Accept} & \text{if } F(x) = y \\ \text{Reject} & \text{if } F(x) \neq y \end{cases}$$

MPC-in-the-Head transform

Zero-knowledge proof



MPCitH transform

Prover

Verifier

MPCitH transform

- ① Generate and commit shares
 $[[x]] = ([[x]]_1, \dots, [[x]]_N)$

$\text{Com}^{\rho_1}([[x]]_1)$
⋮
 $\text{Com}^{\rho_N}([[x]]_N)$

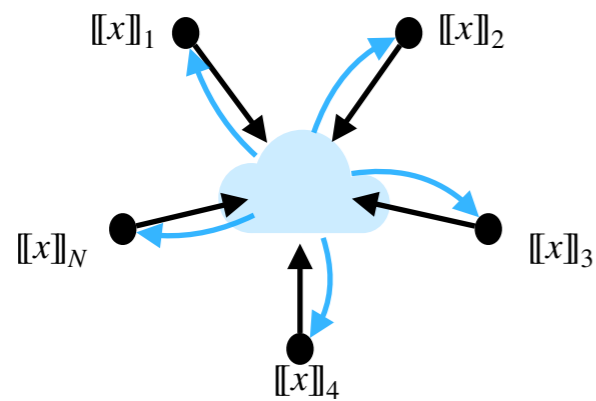
Prover

Verifier

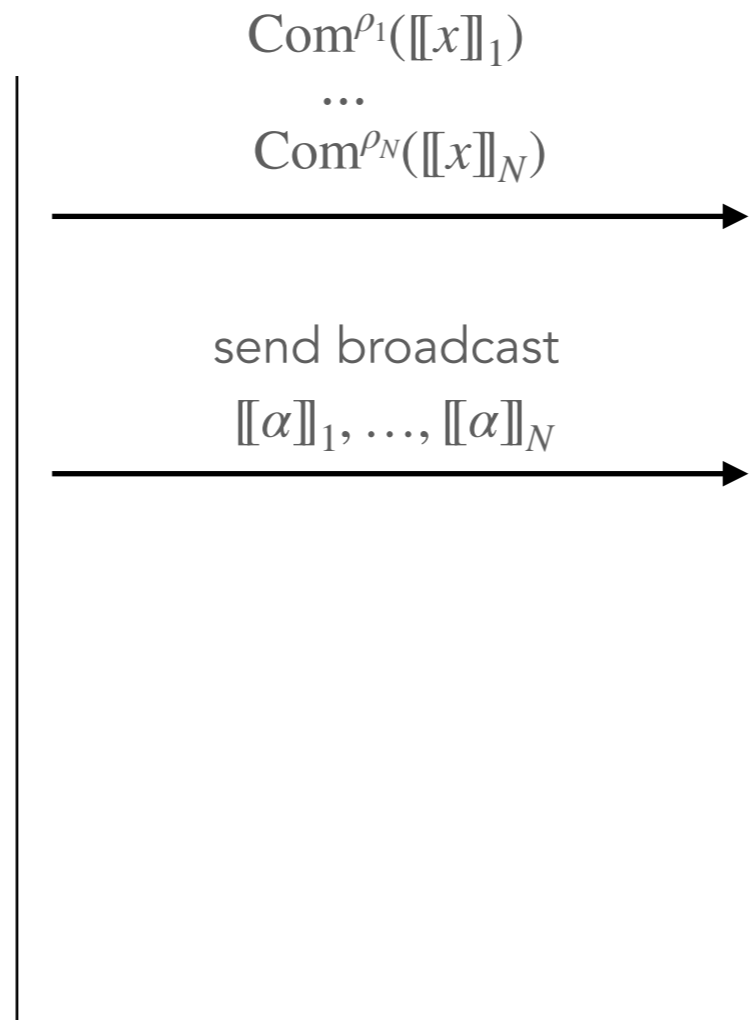
MPCitH transform

- ① Generate and commit shares
 $[[x]] = ([[x]]_1, \dots, [[x]]_N)$

- ② Run MPC in their head



Prover



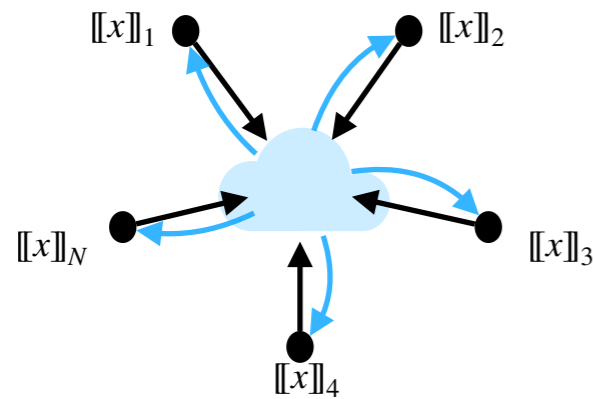
Verifier

MPCitH transform

① Generate and commit shares

$$[[x]] = ([[x]]_1, \dots, [[x]]_N)$$

② Run MPC in their head



Prover

$\text{Com}^{\rho_1}([[x]]_1)$

\dots
 $\text{Com}^{\rho_N}([[x]]_N)$

send broadcast

$[[a]]_1, \dots, [[a]]_N$

i^*

③ Choose a random party

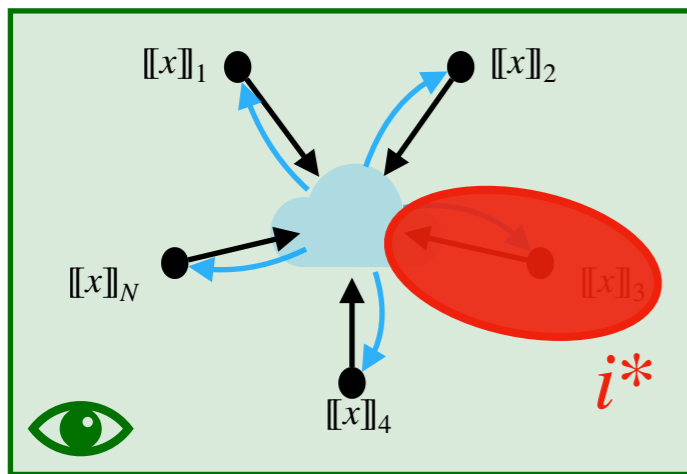
$$i^* \leftarrow^{\$} \{1, \dots, N\}$$

Verifier

MPCitH transform

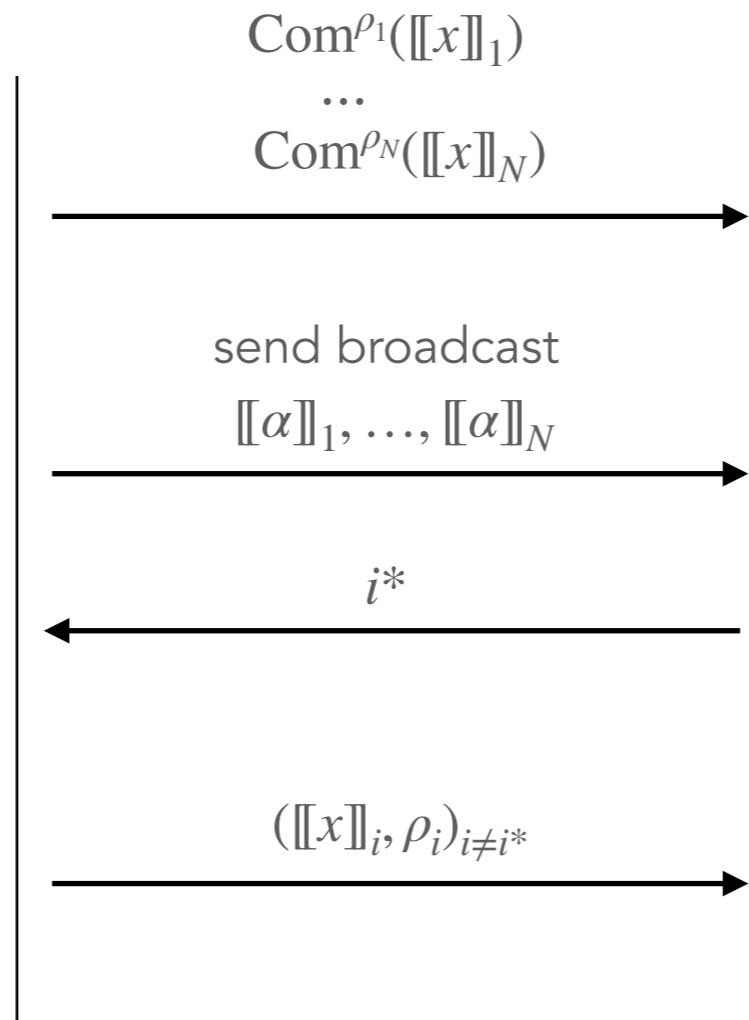
① Generate and commit shares
 $[[x]] = ([[x]]_1, \dots, [[x]]_N)$

② Run MPC in their head



④ Open parties $\{1, \dots, N\} \setminus \{i^*\}$

Prover



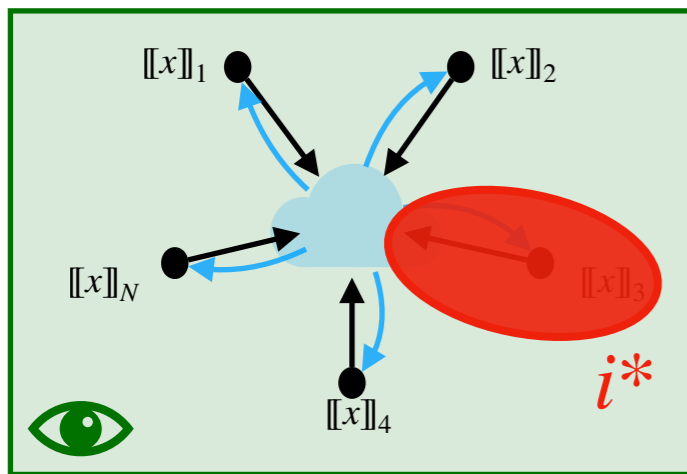
③ Choose a random party
 $i^* \leftarrow^{\$} \{1, \dots, N\}$

Verifier

MPCitH transform

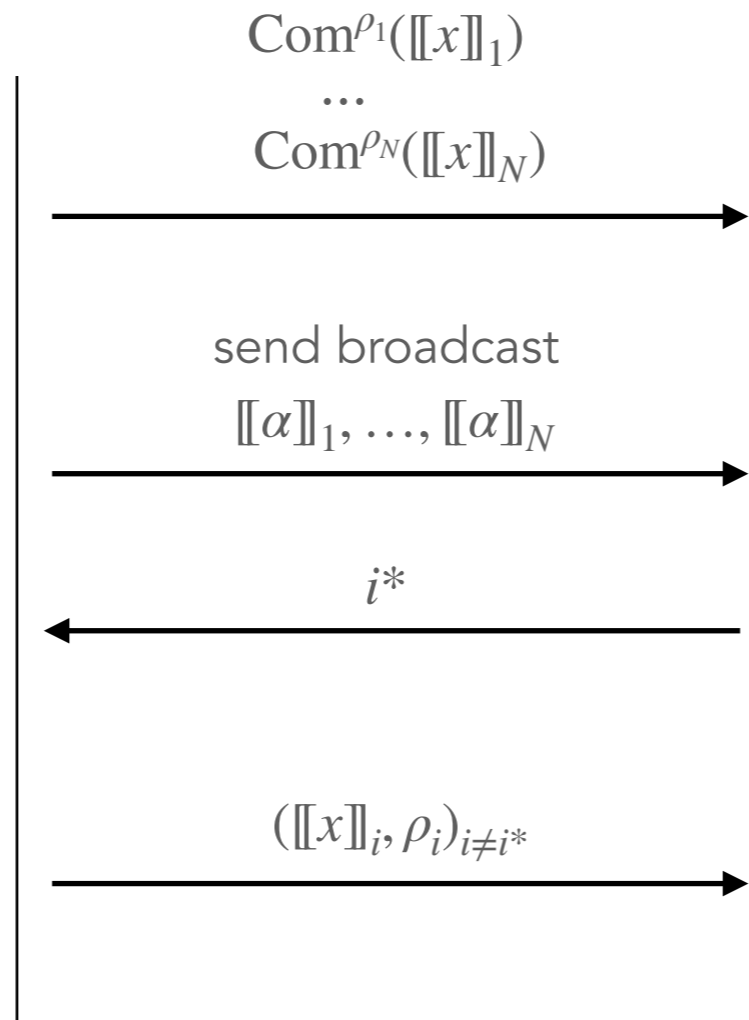
① Generate and commit shares
 $[[x]] = ([[x]]_1, \dots, [[x]]_N)$

② Run MPC in their head



④ Open parties $\{1, \dots, N\} \setminus \{i^*\}$

Prover



③ Choose a random party
 $i^* \leftarrow^{\$} \{1, \dots, N\}$

⑤ Check $\forall i \neq i^*$
 - Commitments $\text{Com}^{\rho_i}([[x]]_i)$
 - MPC computation $[[\alpha]]_i = \varphi([[x]]_i)$
 Check $g(y, \alpha) = \text{Accept}$

Verifier

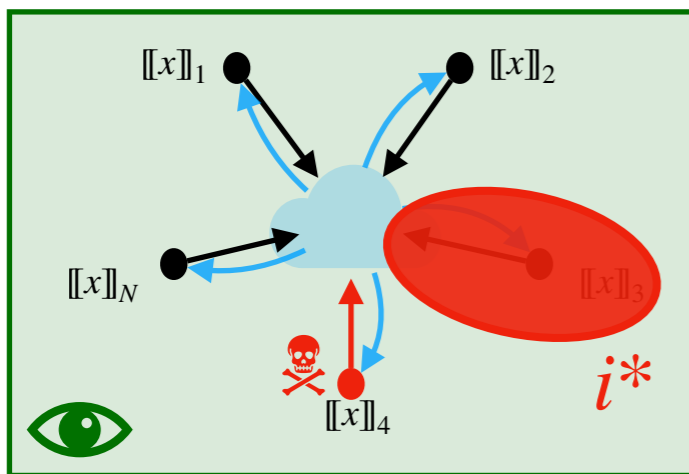
MPCitH transform

① Generate and commit shares

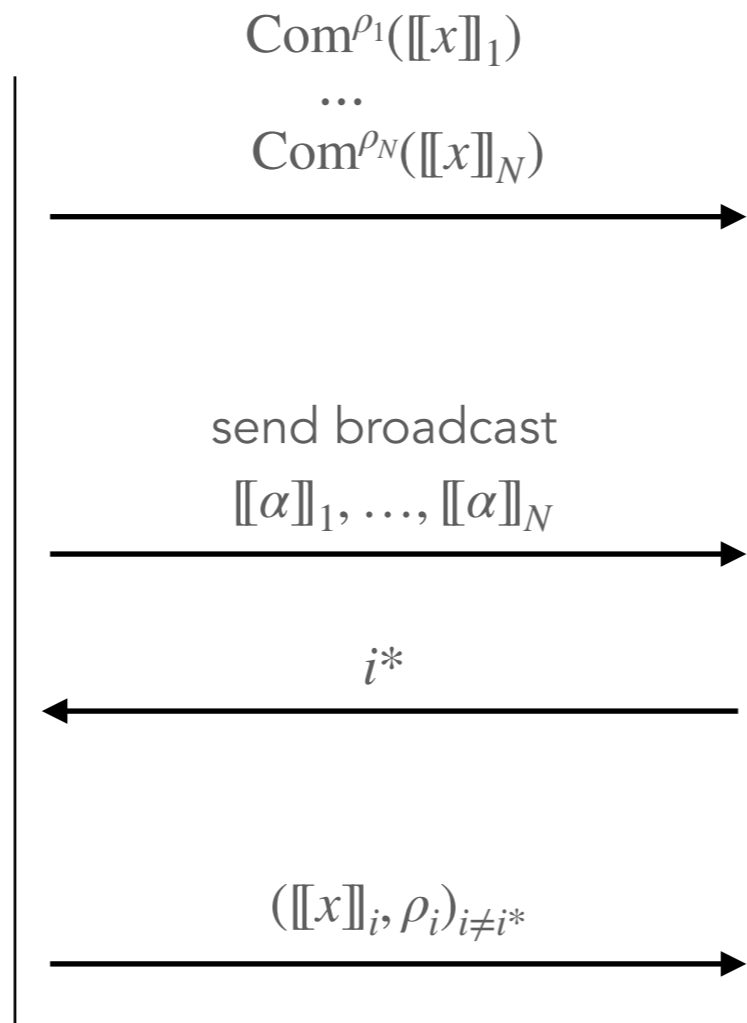
$$[[x]] = ([[x]]_1, \dots, [[x]]_N)$$

We have $F(x) \neq y$ where
 $x := [[x]]_1 + \dots + [[x]]_N$

② Run MPC in their head



④ Open parties $\{1, \dots, N\} \setminus \{i^*\}$



③ Choose a random party
 $i^* \leftarrow^{\$} \{1, \dots, N\}$

⑤ Check $\forall i \neq i^*$
 - Commitments $\text{Com}^{\rho_i}([[x]]_i)$
 - MPC computation $[[\alpha]]_i = \varphi([[x]]_i)$
 Check $g(y, \alpha) = \text{Accept}$

Malicious Prover

Verifier

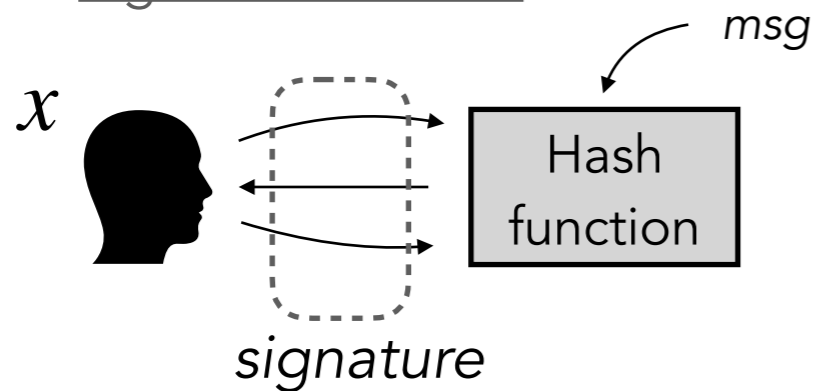
✗ Cheating detected!

One-way function

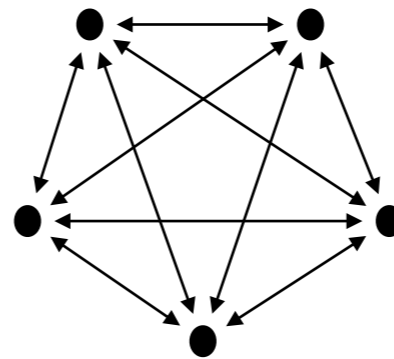
$$F : x \mapsto y$$

E.g. AES, MQ system,
Syndrome decoding

Signature scheme



Multiparty computation (MPC)

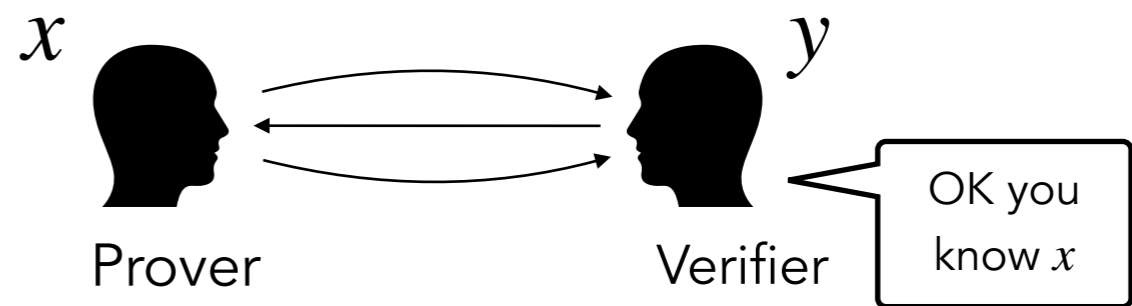


Input sharing $[[x]]$
Joint evaluation of:

$$g(x) = \begin{cases} \text{Accept} & \text{if } F(x) = y \\ \text{Reject} & \text{if } F(x) \neq y \end{cases}$$

MPC-in-the Head transform

Zero-knowledge proof

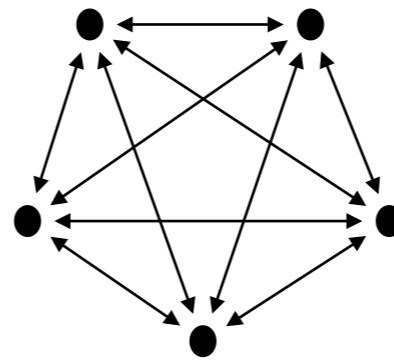


One-way function

$$F : x \mapsto y$$

E.g. AES, MQ system,
Syndrome decoding

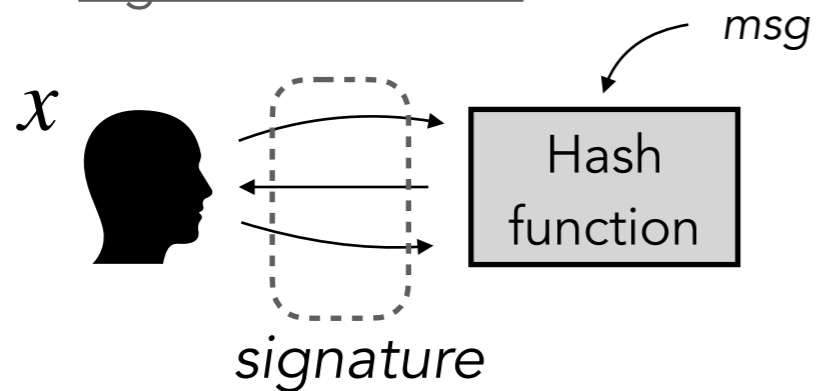
Multiparty computation (MPC)



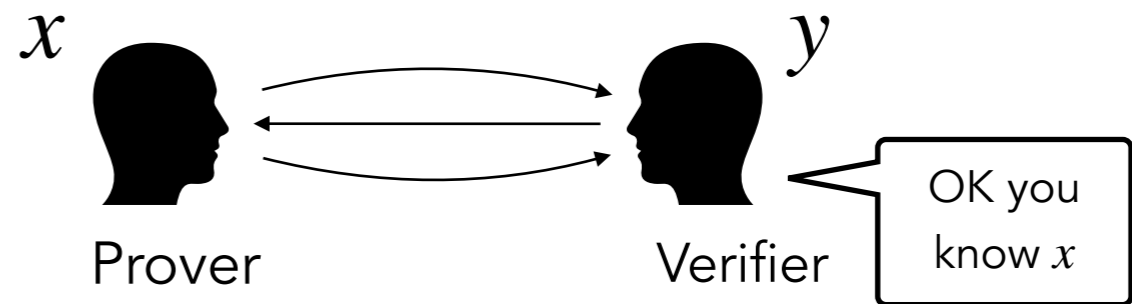
Input sharing $[[x]]$
Joint evaluation of:

$$g(x) = \begin{cases} \text{Accept} & \text{if } F(x) = y \\ \text{Reject} & \text{if } F(x) \neq y \end{cases}$$

Signature scheme



Zero-knowledge proof



Fiat-Shamir transform

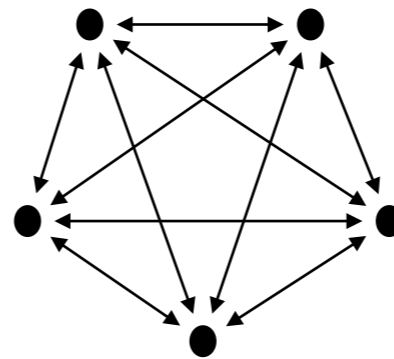


One-way function

$$F : x \mapsto y$$

E.g. AES, MQ system,
Syndrome decoding

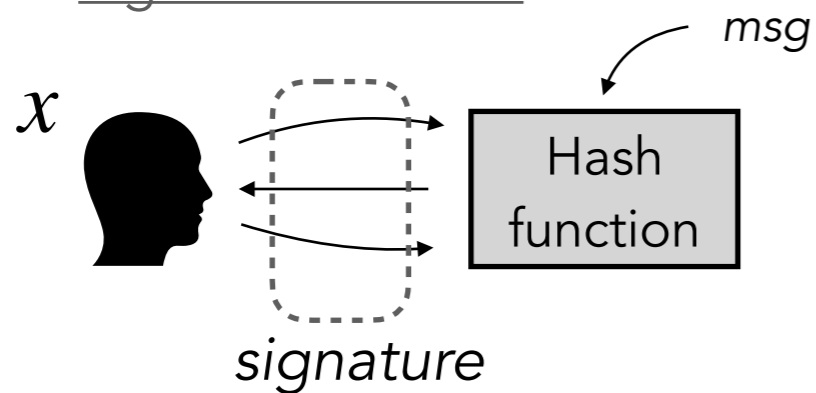
Multiparty computation (MPC)



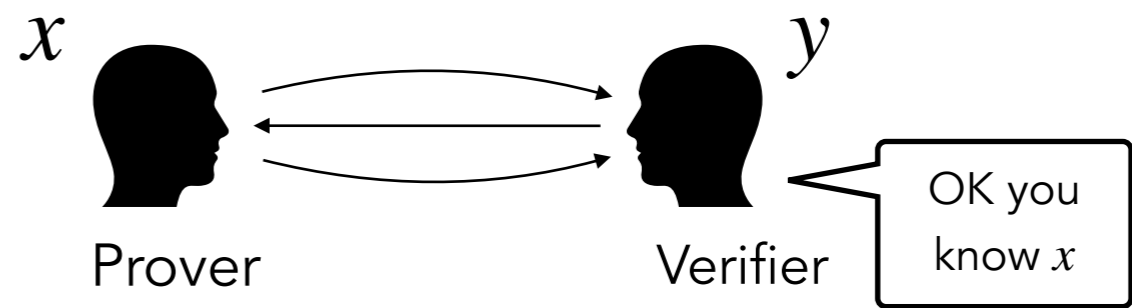
Input sharing $[[x]]$
Joint evaluation of:

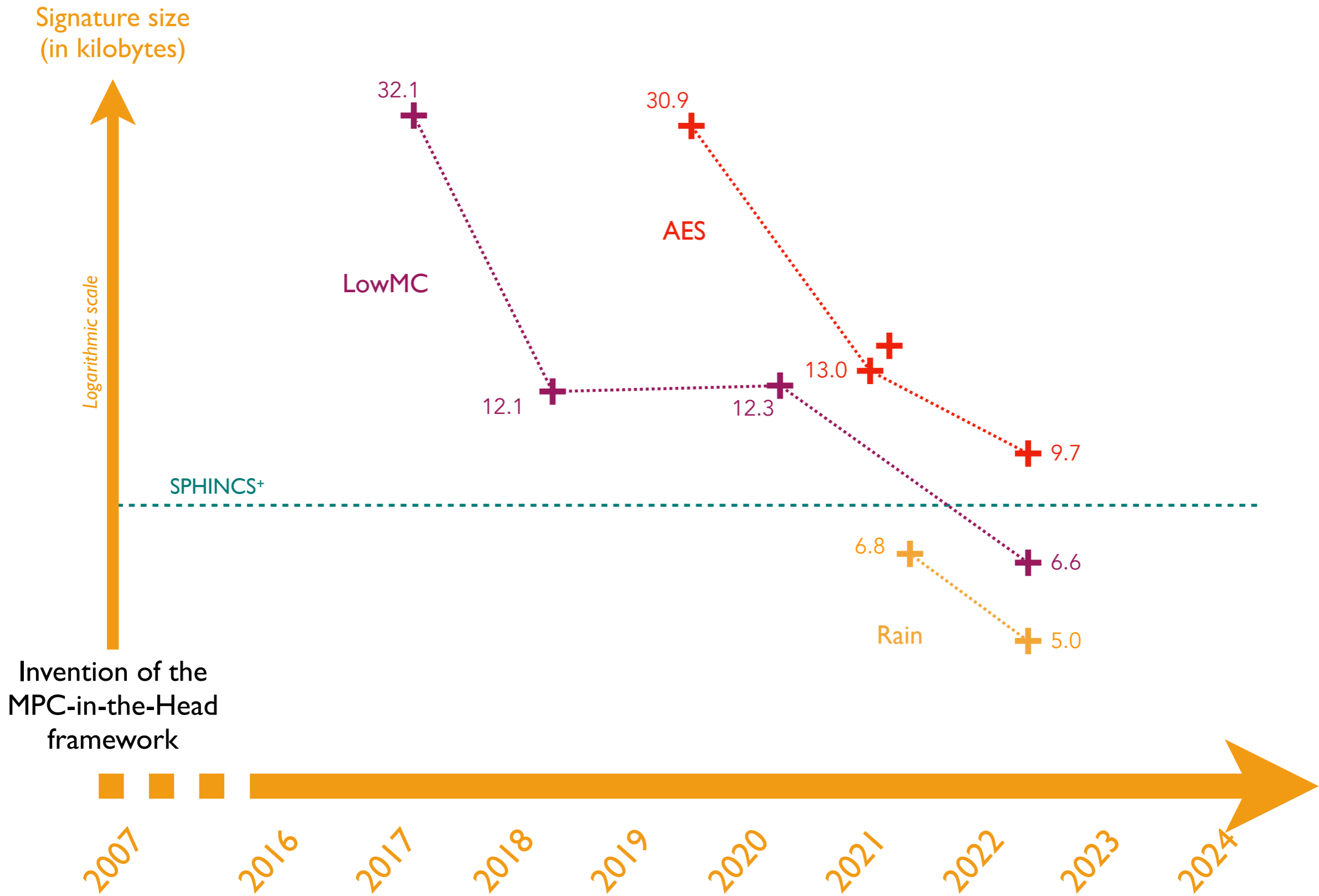
$$g(x) = \begin{cases} \text{Accept} & \text{if } F(x) = y \\ \text{Reject} & \text{if } F(x) \neq y \end{cases}$$

Signature scheme



Zero-knowledge proof





Signature size
(in kilobytes)

Logarithmic scale

SPHINCS+

Syndrome Decoding Problem:

From a matrix H and a vector y , find x such that

- $y = Hx$,
- x has at most w non-zero coordinates.

1990

1995

2000

2005

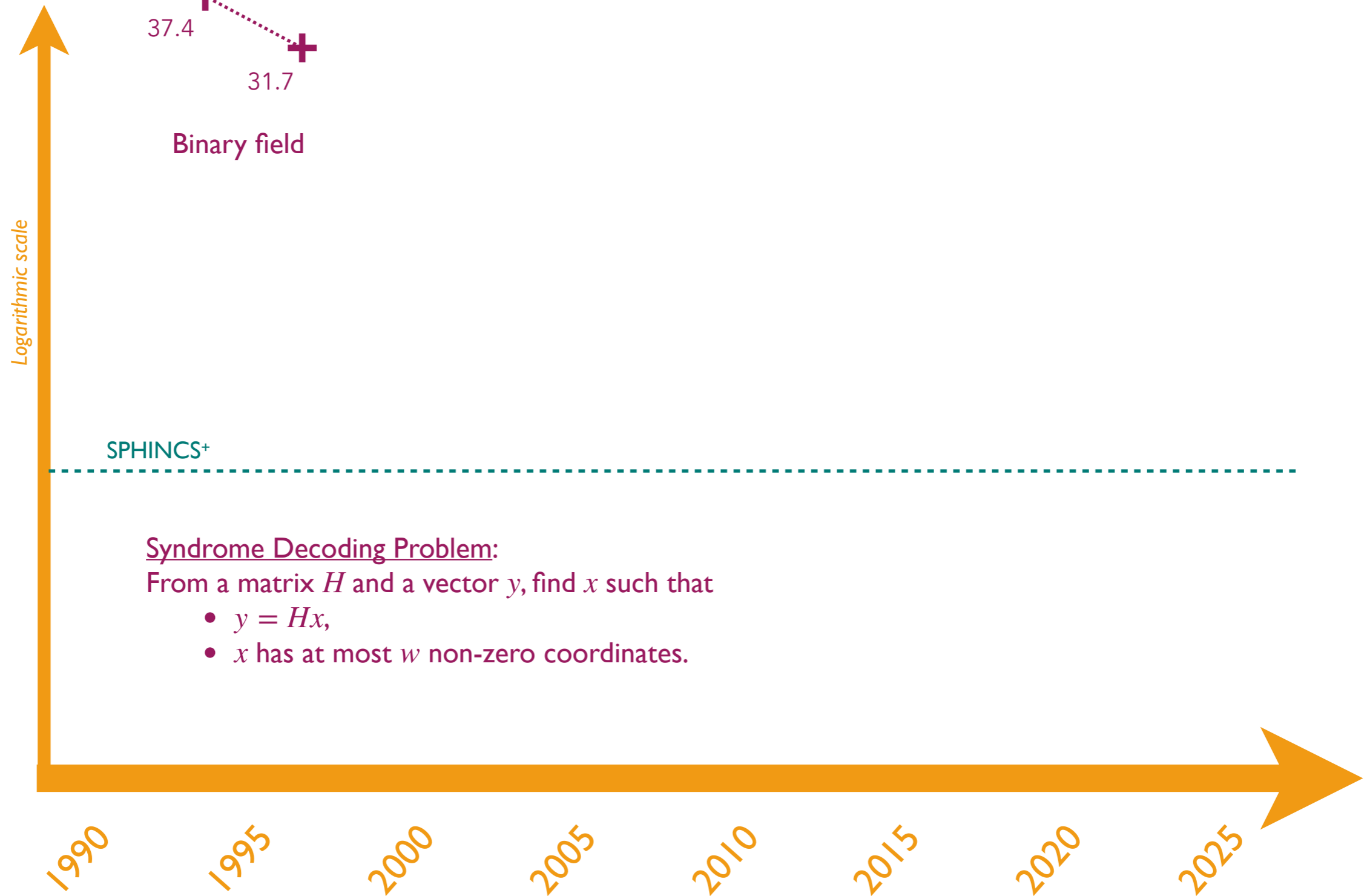
2010

2015

2020

2025

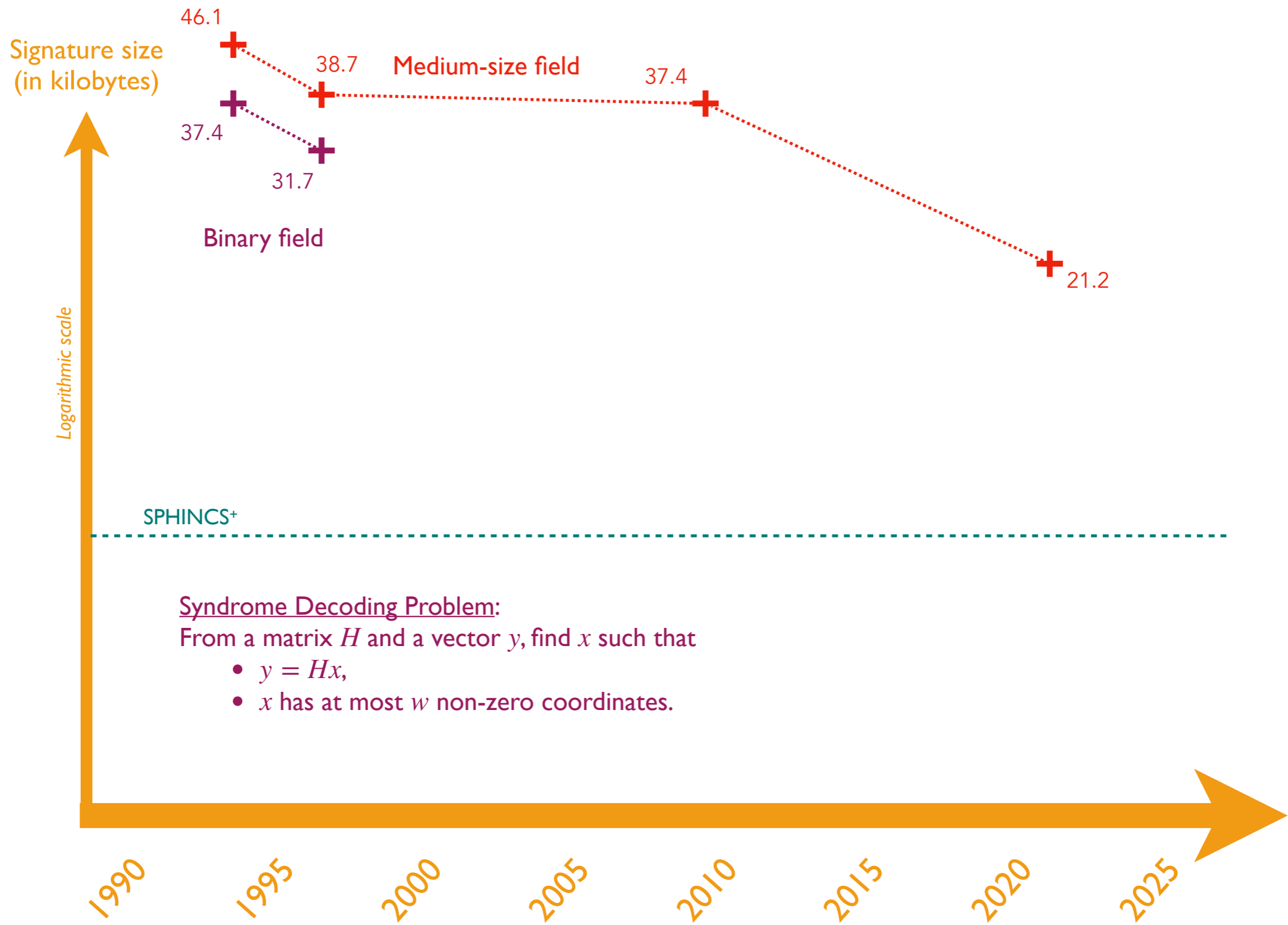
Signature size
(in kilobytes)



Syndrome Decoding Problem:

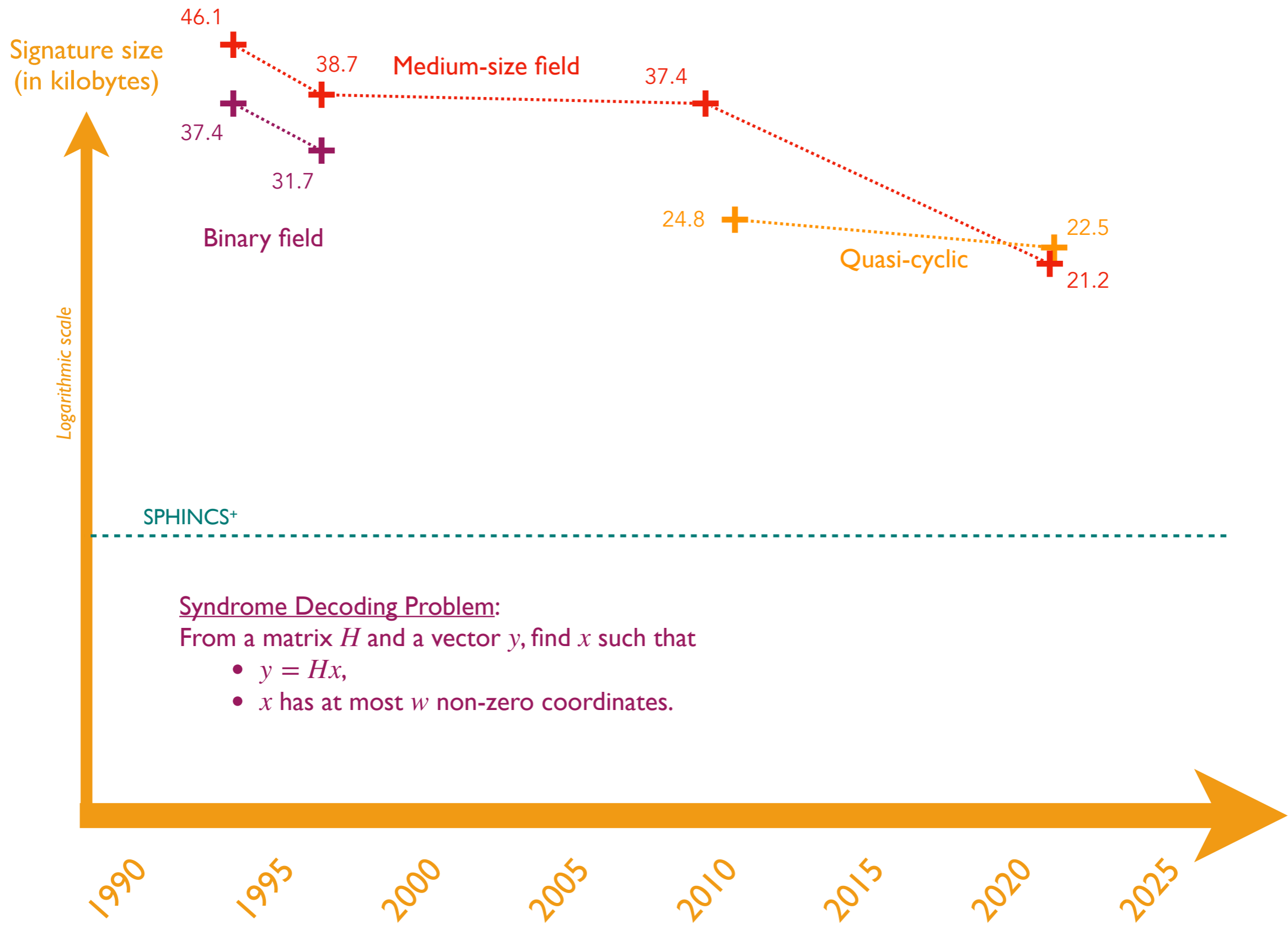
From a matrix H and a vector y , find x such that

- $y = Hx$,
- x has at most w non-zero coordinates.



Syndrome Decoding Problem:
 From a matrix H and a vector y , find x such that

- $y = Hx$,
- x has at most w non-zero coordinates.

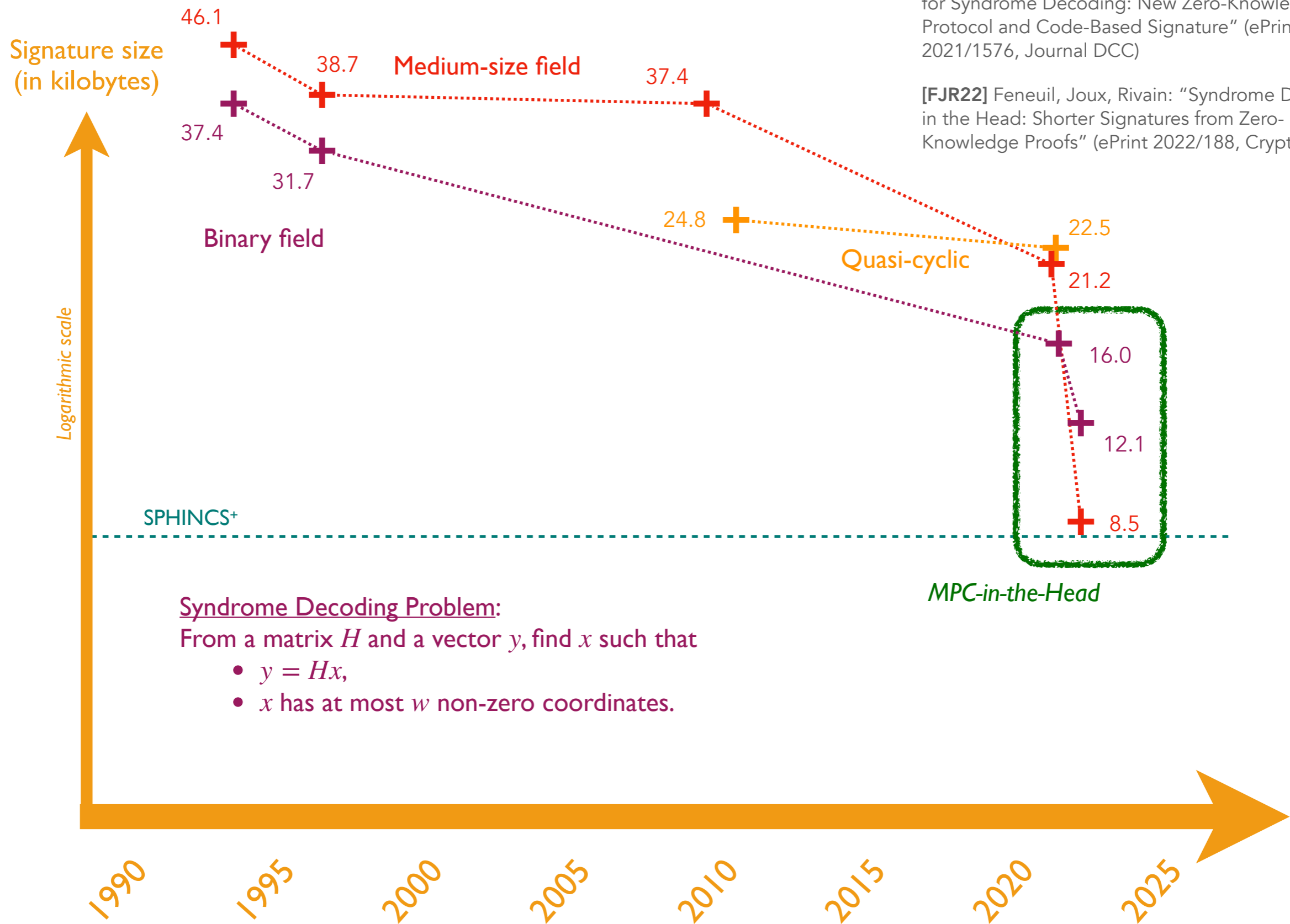


Syndrome Decoding Problem:
 From a matrix H and a vector y , find x such that

- $y = Hx$,
- x has at most w non-zero coordinates.

[FJR23] Feneuil, Joux, Rivain: "Shared Permutation for Syndrome Decoding: New Zero-Knowledge Protocol and Code-Based Signature" (ePrint 2021/1576, Journal DCC)

[FJR22] Feneuil, Joux, Rivain: "Syndrome Decoding in the Head: Shorter Signatures from Zero-Knowledge Proofs" (ePrint 2022/188, Crypto 2022)



Syndrome Decoding Problem:
From a matrix H and a vector y , find x such that

- $y = Hx$,
- x has at most w non-zero coordinates.

Exploring other assumptions

- Subset Sum Problem: ≥ 100 KB \Rightarrow 19.1 KB
- Multivariate Quadratic Problem: 6.3 – 7.3 KB
- MinRank Problem: $\approx 5 - 6$ KB
- Rank Syndrome Decoding Problem: $\approx 5 - 6$ KB
- Permuted Kernel Problem (or variant): ≈ 6 KB
- ...

MPCitH-based NIST Candidates

1st June 2023:

Deadline for the NIST call
for additional post-quantum signatures

MPCitH-based NIST Candidates

	Assumption	Size (in KB)
AIMer	AIM (MPC-friendly one-way function)	4.2
Biscuit	Structured MQ problem (PowAff2)	4.7
MIRA	MinRank problem	5.6
MiRitH	MinRank problem	5.7
RYDE	Syndrome decoding problem in rank metric	6.0
PERK*	Permuted Kernel problem (variant)	6.1
MQOM	Unstructured MQ problem	6.3
SDitH	Syndrome decoding problem in Hamming	8.2

MPCitH-based NIST Candidates

	Assumption	Size (in KB)
AIMer	AIM (MPC-friendly one-way function)	4.2
Biscuit	Structured MQ problem (PowAff2)	4.7
MIRA	MinRank problem	5.6
MiRitH	MinRank problem	5.7
RYDE	Syndrome decoding problem in rank metric	6.0
PERK	Permuted Kernel problem (variant)	6.1
MQOM	Unstructured MQ problem	6.3
SDitH	Syndrome decoding problem in Hamming	8.2



What about the computational cost ?

Traditional Transformation

(2018) Emulation : N parties

Traditional Transformation

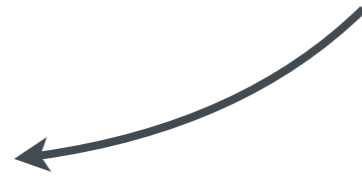
(2018) Emulation : N parties

Hypercube Technique

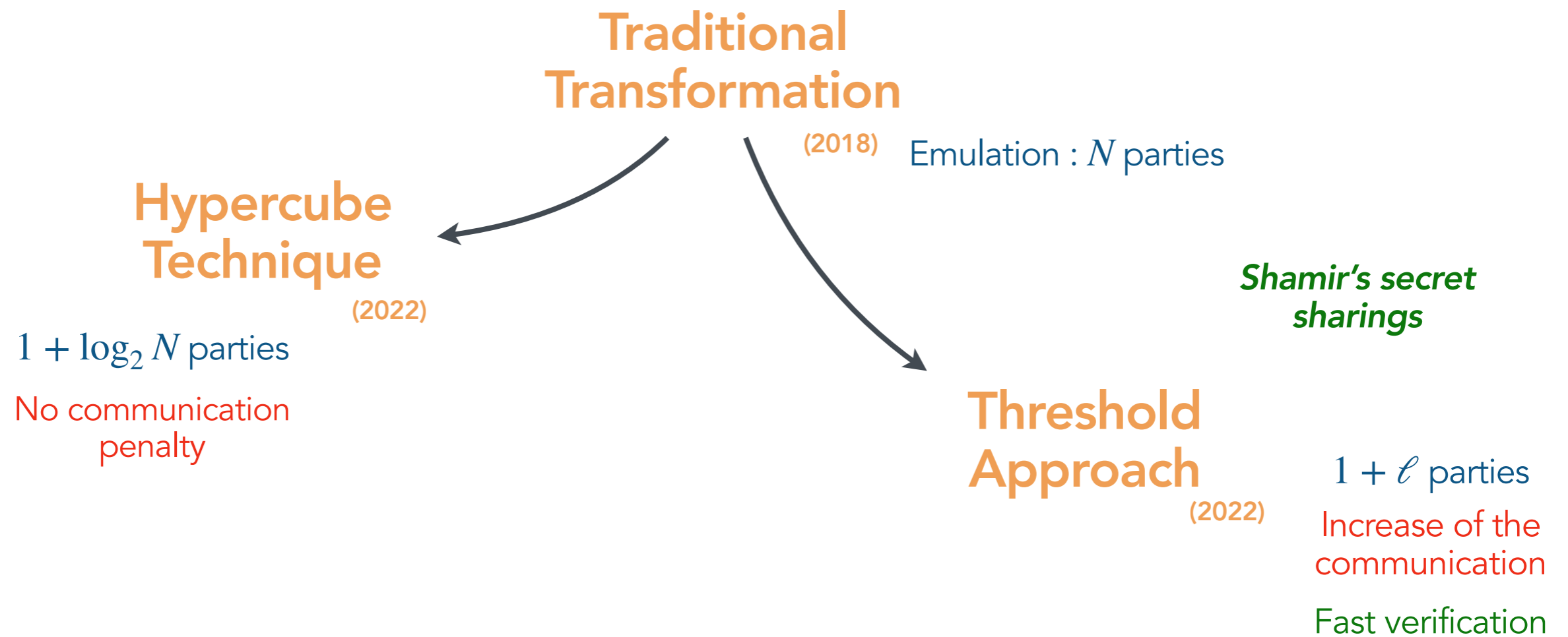
(2022)

$1 + \log_2 N$ parties

No communication penalty



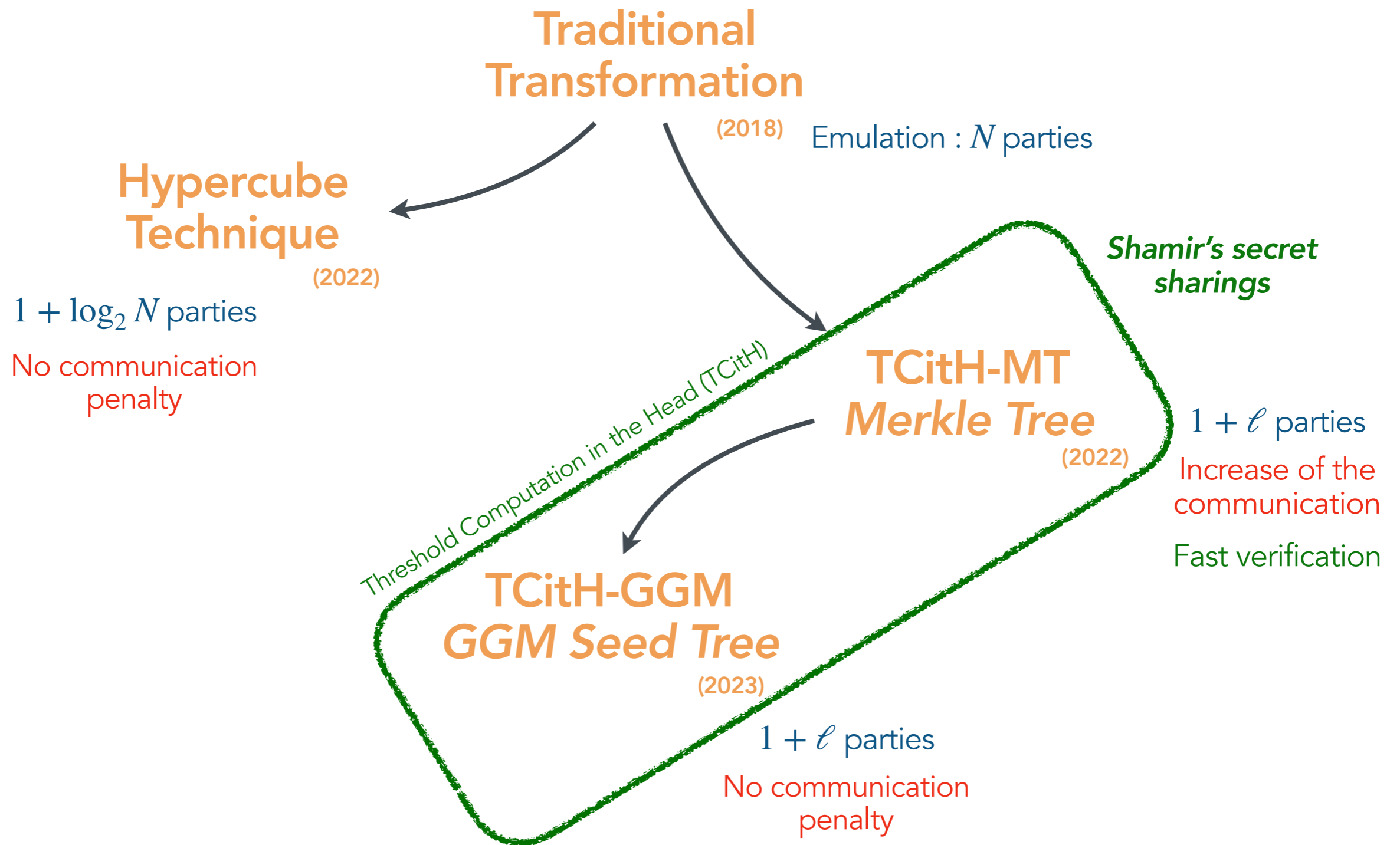
[AGHHJY23] Aguilar-Melchor, Gama, Howe, Hülsing, Joseph, Yue: "The Return of the SDitH" (ePrint 2022/1645, Eurocrypt 2023)



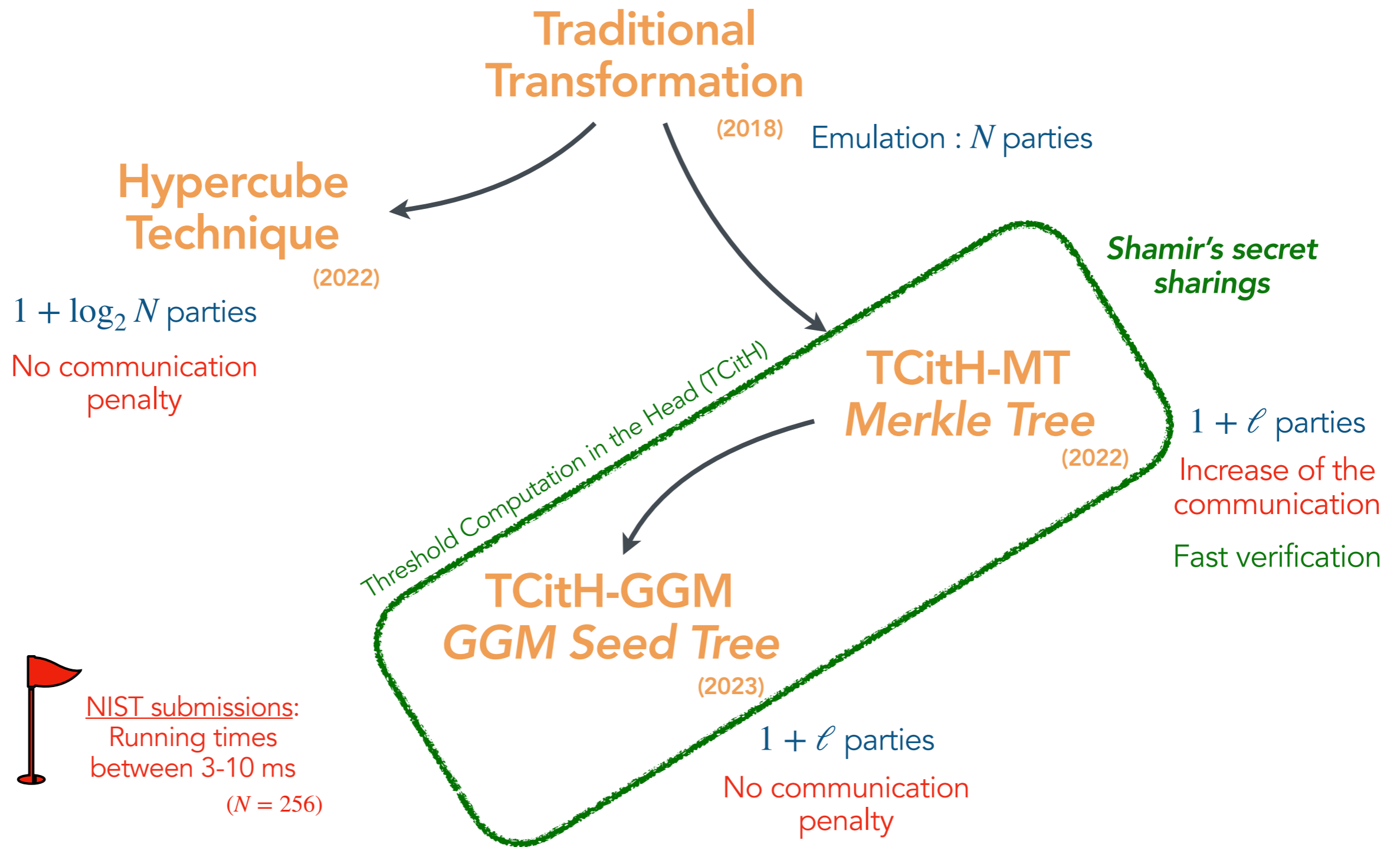
Shamir's secret sharing: to share a value s ,

- Build a random degree- ℓ polynomial $P(X) := s + \sum_{j=1}^{\ell} r_j X^j$.
- Set the i^{th} share $[[s]]_i$ as $[[s]]_i := P(e_i)$, where $e_i \neq 0$.

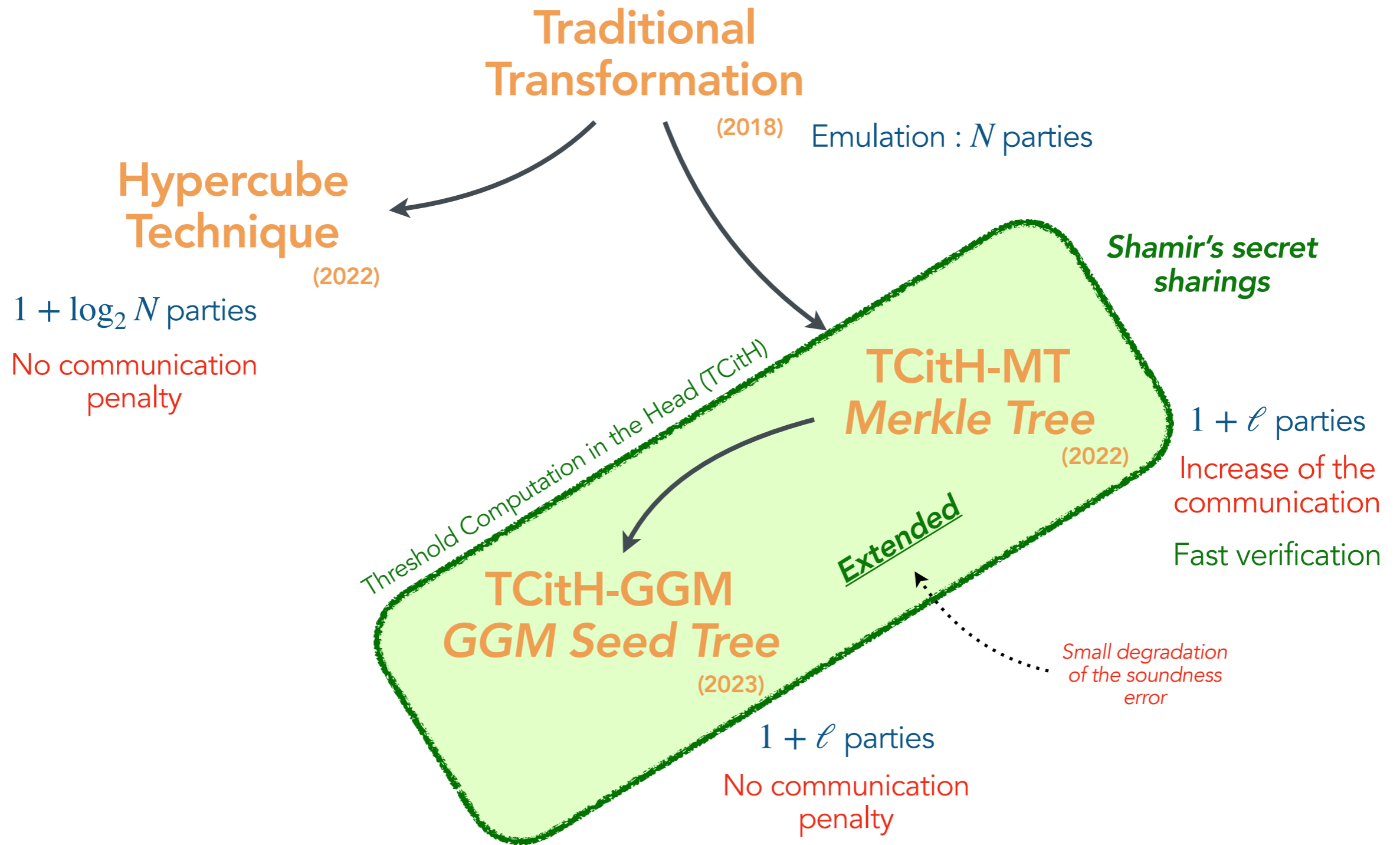
[FR22] Feneuil, Rivain: "Threshold Linear Secret Sharing to the Rescue of MPC-in-the-Head" (ePrint 2022/1407, Asiacrypt 2023)



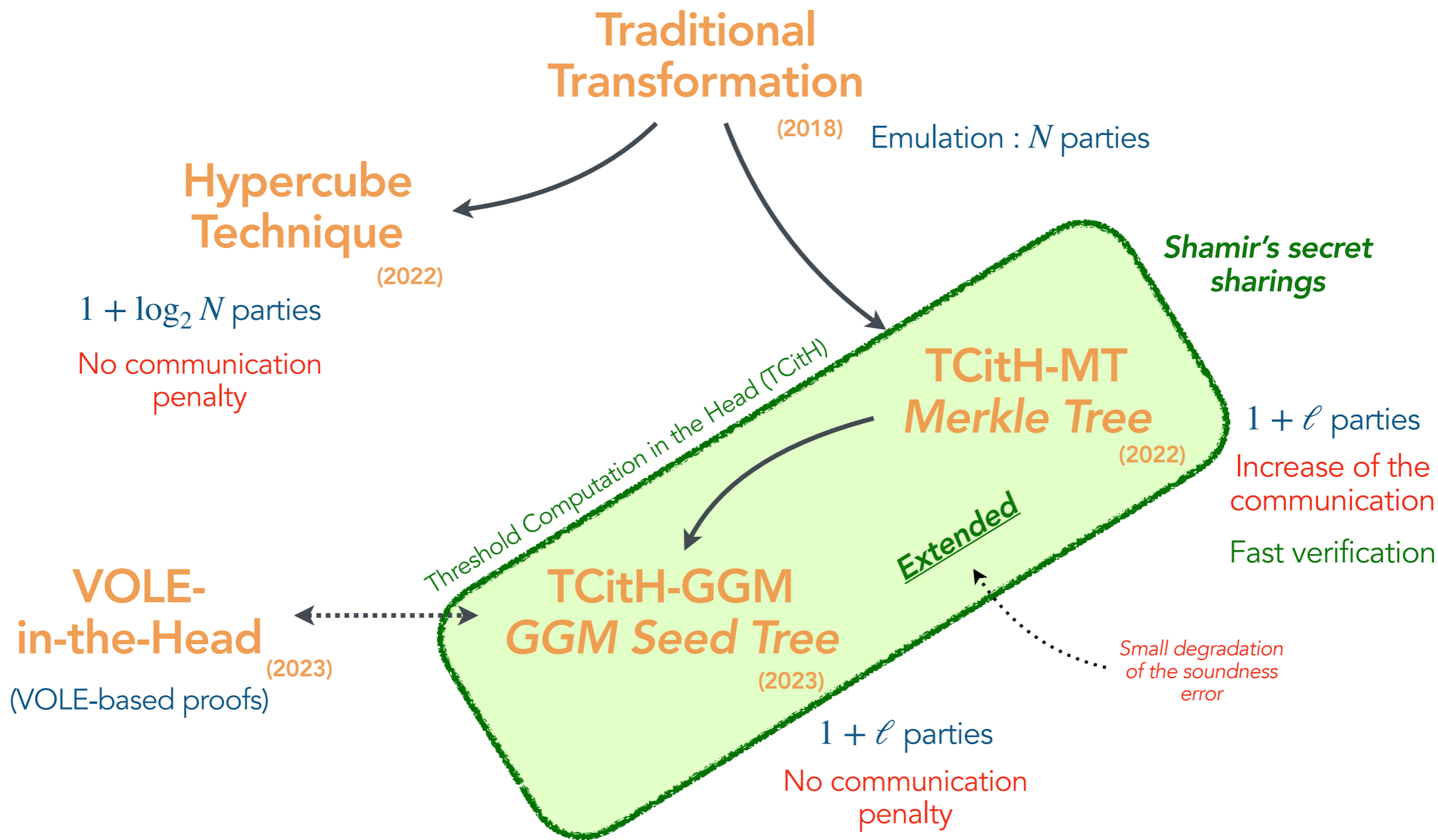
[FR23] Feneuil, Rivain: "Threshold Computation in the Head: Improved Framework for Post-Quantum Signatures and Zero-Knowledge Arguments" (ePrint 2023/1573)



[FR23] Feneuil, Rivain: "Threshold Computation in the Head: Improved Framework for Post-Quantum Signatures and Zero-Knowledge Arguments" (ePrint 2023/1573)



[FR23] Feneuil, Rivain: "Threshold Computation in the Head: Improved Framework for Post-Quantum Signatures and Zero-Knowledge Arguments" (ePrint 2023/1573)



[BBDKORS23] Baum, Braun, Delpech, Klooß, Orsini, Roy, Scholl: "Publicly Verifiable Zero-Knowledge and Post-Quantum Signatures and VOLE-in-the-Head" (Crypto 2023)

Extended TCitH: some applications

[FR23] Feneuil, Rivain: "Threshold Computation in the Head: Improved Framework for Post-Quantum Signatures and Zero-Knowledge Arguments" (ePrint 2023/1573)

Extended TCitH: some applications

- More efficient signature schemes
 - *Unstructured multivariate quadratic (MQ) problem over \mathbb{F}_{251}*
 - MQOM: 6.5 KB
 - Extended TCitH: 4.2 KB

[FR23] Feneuil, Rivain: "Threshold Computation in the Head: Improved Framework for Post-Quantum Signatures and Zero-Knowledge Arguments" (ePrint 2023/1573)

Extended TCitH: some applications

- More efficient signature schemes
 - *Unstructured multivariate quadratic (MQ) problem over \mathbb{F}_{251}*
 - MQOM: **6.5 KB**
 - Extended TCitH: **4.2 KB**
- Shorter post-quantum ring signature schemes
 - Extended TCitH with MQ: **5.8 KB** in around 8 ms, for 4000 users
 - Extended TCitH with SD: **10.30 KB** in around 10 ms, for 4000 users

[FR23] Feneuil, Rivain: "Threshold Computation in the Head: Improved Framework for Post-Quantum Signatures and Zero-Knowledge Arguments" (ePrint 2023/1573)

Conclusion

■ MPC-in-the-Head

- Very versatile and tunable
- A practical tool to build *conservative* signature schemes
 - *Between 4-10 KB in few milliseconds (NIST Level I)*

■ Perspectives

- Active research field
- Signatures with advanced functionalities:
 - ring signatures, threshold signatures, multi-signatures,
 - blind signatures, ...

PhD Defense

- Title:

Post-Quantum Signatures
from Secure Multiparty Computation

- When: Monday 23rd October 2023, at 2 pm

- Where: Sorbonne University (Jussieu)

More information at



<https://www.thibauld-feneuil.fr/phd-defense.html>